

A numerical introduction to tensor networks for quantum simulation

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[Credit for course materials: Prof. Jan von Delft]

1. Why tensor networks?

Tensor networks provide a flexible description of quantum states.

In some cases, they are efficient - can accurately describe state with polynomial resources.

Example: spin- s chain, with N sites:



Local state space of site l : $|\sigma_l\rangle_l \in \{|\uparrow\rangle_l, |\downarrow\rangle_l, \dots, |2s+1\rangle_l\}$

Local state label: $\sigma = 1, 2, \dots, 2s+1$

Local dimension: $d \equiv 2s+1$

Shorthand: $|\sigma_l\rangle \equiv |\sigma_l\rangle_l$

Index l on state label σ_l suffices to identify the site Hilbert space $|\uparrow\rangle$

Generic basis state for full chain of length N
(convention: add state spaces for new sites from left):

$$|\sigma_N\rangle \otimes \dots \otimes |\sigma_2\rangle \otimes |\sigma_1\rangle \equiv |\sigma_1 \sigma_2 \dots \sigma_N\rangle \equiv |\vec{\sigma}\rangle_N$$

Hilbert space for full chain:

$$\mathcal{H}^N = \text{span}(\{|\vec{\sigma}\rangle_N\}) \quad \leftarrow \text{superpositions}$$

Generic quantum state:

$$(|\psi\rangle_N \in \mathcal{H}^N) \quad |\psi\rangle_N = \sum_{\sigma_1 \dots \sigma_N} |\sigma_1 \dots \sigma_N\rangle C^{\sigma_1 \dots \sigma_N} \equiv \underbrace{C^{\vec{\sigma}}}_{\text{summation implied}} |\vec{\sigma}\rangle_N$$

Dimension of full Hilbert space

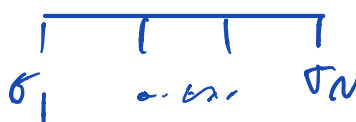
$$\mathcal{H}^N : \underbrace{d^N}_{\text{exponential explosion}} \quad (\# \text{ of different configurations of } \vec{\sigma})$$

Specifying $|\psi\rangle_N$ involves specifying $C^{\vec{\sigma}}$, i.e. d^N different complex numbers.

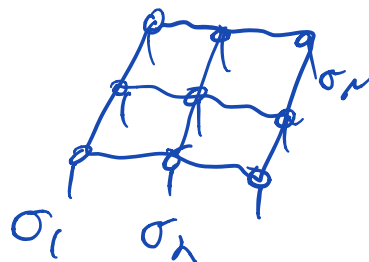
$C^{\vec{\sigma}} = C^{\sigma_1 \dots \sigma_N}$ is a tensor of rank N (rank = # of legs)

Graphical representation: (one leg for each index)

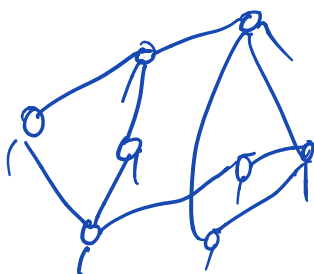
$$C^{\vec{\sigma}} \equiv \text{diagram of a tensor with } N \text{ legs labeled } \sigma_1, \sigma_2, \dots, \sigma_N$$

Claim: such a rank L tensor can be represented in many different ways:

MPS: matrix product state



PEPS: projected entangled-pair state



arbitrary tensor network

- a link between two sites represents entanglement between them
- different representations \rightarrow different entanglement book-keeping
- tensor network = entanglement representation of a quantum state

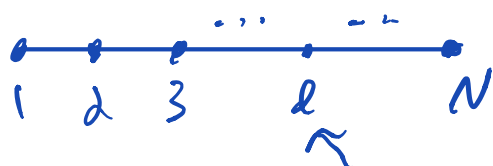
virtual indices

2. Iterative diagonalization

Consider a spin- s chain with Hamiltonian

$$H^N = \sum_{l=1}^{N-1} \vec{S}_l \cdot \vec{S}_{l+1} + \sum_{l=1}^N \vec{S}_l \cdot \vec{h}_l$$

(example)



local state space for site l : $|\sigma_l\rangle = |1, 2, \dots, d\rangle$ $d = 2s+1$

We seek eigenstates of H^N : $H^N |E_\alpha^N\rangle = E_\alpha^N |E_\alpha^N\rangle$

$|E_\alpha^N\rangle$ = eigenvector of H^N

E_α^N = eigenvalue

$|E_\alpha^N\rangle \in \mathcal{H}^N$; $\alpha = 1 \dots d^N$

Diagonalize Hamiltonian iteratively, adding one site at a time.

N = 1: Start with first site, diagonalize H^1 in Hilbert space \mathcal{H}^1 . Eigenstates have form

$$|\alpha\rangle \equiv |E'_\alpha\rangle = |\sigma_1\rangle A_{\alpha}^{\sigma_1} \quad (\alpha = 1, \dots, d)$$

\uparrow \uparrow \uparrow
 (sum over σ_1 implied) coefficient matrix A combining 'incoming' σ_1 into 'outgoing' α

N = 2: Add second site, diagonalize H^2 in Hilbert space \mathcal{H}^2 :

$$|\beta\rangle \equiv |E'_\beta\rangle = |\sigma_2\rangle \otimes |\alpha\rangle B_{\beta}^{\alpha\sigma_2} \quad (\beta = 1, \dots, d^2)$$

(sum over α, σ_2 implied) coefficient matrix ~~tensor~~ combining 'incoming' α, σ_2 into 'outgoing' β

$$|\beta\rangle = |\sigma_2\rangle \otimes |\sigma_1\rangle A_{\alpha}^{\sigma_1} B_{\beta}^{\alpha\sigma_2}$$

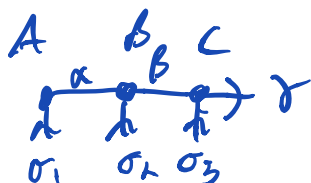
$\underbrace{\hspace{10em}}$
 contraction \rightarrow matrix multiplication

N = 3: Add third site, diagonalize H^3 in Hilbert space \mathcal{H}^3 :

Your try:

Answer: $|\gamma\rangle = |\sigma_3\rangle \oplus |\beta\rangle \subset^{\beta\sigma_3} \gamma \quad (\gamma = 1 \dots d^3)$

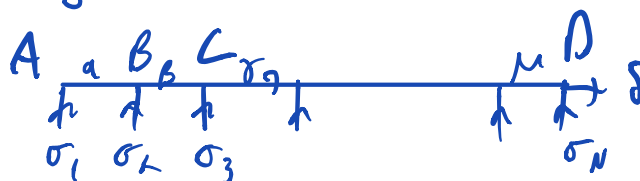
$$= |\sigma_3\rangle \oplus |\sigma_2\rangle \oplus |\sigma_1\rangle \underbrace{A^{\sigma_1}_\alpha B^{\alpha\sigma_2}_\beta \subset^{\beta\sigma_3}}_{\text{Contracted}} \gamma$$



Continue similarly until having added site N. Eigenstates of H^N have following structure:

$$|E^N_\delta\rangle = |\delta\rangle = |\sigma_N\rangle \oplus \dots \oplus |\sigma_2\rangle \oplus |\sigma_1\rangle \underbrace{A^{\sigma_1}_\alpha B^{\alpha\sigma_2}_\beta \dots D^{\mu\sigma_N}_\delta}_{\subset^{\bar{\sigma}}_\delta} \\ = |\bar{\sigma}_N\rangle \subset^{\bar{\sigma}}_\delta \quad (\text{MPS})$$

$$(\delta = 1 \dots d^N)$$



Nomenclature: σ_ℓ = physical indices, $\alpha, \beta, \gamma, \dots$ = (virtual) bond indices

Alternative, widely-used notation: 'reshape' coefficient tensors as

$$\tilde{A}^{\sigma_1}_\alpha \equiv A^{\sigma_1}_\alpha \quad ; \quad \tilde{B}^{\sigma_2}_{\alpha\beta} \equiv B^{\alpha\sigma_2}_\beta$$

to highlight 'matrix product' structure in noncovariant notation:

$$|\delta\rangle = |\sigma_N\rangle \oplus \dots \oplus |\sigma_1\rangle \underbrace{\tilde{A}^{\sigma_1}_\alpha \tilde{B}^{\sigma_2}_{\alpha\beta} \tilde{C}^{\sigma_3}_{\beta\gamma}} \dots \tilde{D}^{\sigma_N}_{\mu\delta}$$

Comments

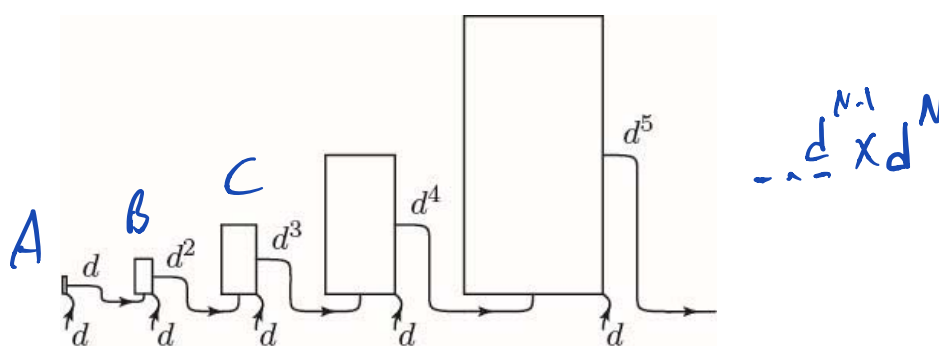
1. Iterative diagonalization of 1D chain generates eigenstates whose wave functions are tensors that are expressed as matrix products \rightarrow matrix product states (MPS)

Matrix size grows exponentially.

for given σ_1 , $A_{\alpha}^{\sigma_1}$ has dimension $1 \times d$ (vector)

for given σ_2 , $B_{\beta}^{\alpha\sigma_2}$ has dimension $d \times d^2$ (rectangular matrix)

for given σ_3 , $C_{\delta}^{\beta\sigma_3}$ has dimension $d^2 \times d^3$ (larger rectangular matrix)

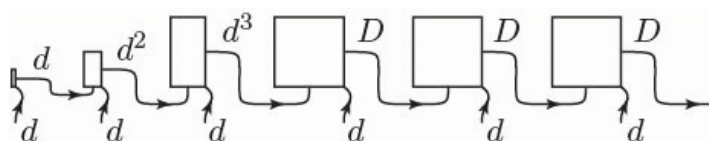


“Hilbert space is a large place!”

Numerical costs explode with increasing N , so truncation schemes are needed.

Truncation can be done in a controlled way using tensor network methods.

Standard truncation scheme: use $\alpha, \beta, \gamma \dots \leq D$ for all virtual bonds



2. Number of parameters available to encode state:

$$N_{\text{MPS}} \leq N \cdot D^2 \cdot d$$

$\hat{L} =$ if all virtual bonds
have dim. D

$$A_{i_1 \dots i_N}^{j_1 \dots j_N}$$

[equal if all virtual bonds have same dimension D]

If N is large: $N_{\text{MPS}} \ll d^N$

Why should we expect this ansatz with polynomial parameters to accurately represent a quantum state?

$$\underline{E_1 - E_0} > \delta \text{ as } N \rightarrow \infty$$

Remarkable fact: for 1d Hamiltonian with local interactions and a gapped spectrum, its ground state can be accurately represented by MPS!

Reason: Area law for entanglement entropy. We will discuss later.

3. Covariant index notation

More detail of covariant index notation is in L2 & L10 of "Mathematics for Physicists", Altland & von Delft, see [here](#).

Kets (Hilbert space vectors)

For kets, indices are down. E.g. basis kets:

Convention

$$|\phi_\sigma\rangle$$

↑
down

For components of kets (wrt a basis), indices are upper:

$$|\psi\rangle = |\phi_\sigma\rangle A^\sigma$$

Repeated indices (always up-down pairs) are always summed (implied summation).

$$\sum_\sigma$$

Example: linear combination of kets.

Your try:

Answer:

$$|\psi_\alpha\rangle = |\phi_\sigma\rangle A^\sigma_\alpha$$

↑
produce a ket

Note: for A^σ_α the index σ identifies components of kets \rightarrow upper
the index α identifies components of basis kets \rightarrow lower

Basis for direct product space: $|\phi_{\vec{\sigma}}\rangle \equiv |\phi_{\sigma_1 \dots \sigma_N}\rangle \equiv |\phi_N\rangle \otimes \dots \otimes |\phi_1\rangle$

Note ket order: start with first space on very right, successively attach new spaces from left.

Linear combinations: $|\psi_{\beta}\rangle = |\phi_{\sigma_1 \dots \sigma_N}\rangle A^{\sigma_1 \dots \sigma_N}_{\beta}$

Your try:

$$= |\phi_{\vec{\sigma}}\rangle A^{\vec{\sigma}}_{\beta}$$

Answer:

Bras (Hilbert space dual vectors)

For bras, indices are upper. E.g. basis bras:

$$\langle \phi^{\sigma} |$$

For components of bras (wrt a basis), indices are lower:

$$\langle \psi | = A^{\dagger}_{\sigma} \langle \phi^{\sigma} |$$

Complex conjugation:

$$A^{\dagger}_{\sigma} = \overline{A^{\sigma}} \quad \leftarrow \text{complex conjugation}$$

Linear combinations of bras:

$$\langle \psi^\alpha | = A_{\sigma}^{\dagger \alpha} \langle \phi^{\sigma} |$$

Your try:

Answer:

Complex conjugation:

$$A_{\sigma}^{\dagger \alpha} = \overline{A_{\alpha}^{\sigma}} \quad (\text{Hermitian conjugate})$$

Note: for $A_{\sigma}^{\dagger \alpha}$, the index α identifies basis bras (dual vectors) hence upper (out)
the index σ identifies components of bras (dual vectors) hence lower

Basis for direct product space:

$$\langle \phi^{\sigma} | \equiv \langle \phi^{\sigma_1} \dots \sigma_N | \equiv \langle \phi^{\sigma_1} | \otimes \dots \otimes \langle \phi^{\sigma_N} |$$

opposite order

Note bra order: opposite to kets so expectation values yield nested bra-ket pairs:

$$\begin{aligned} & \langle \phi^{\sigma_1} \dots \sigma_N | \hat{O} | \phi^{\sigma_1} \dots \sigma_N \rangle \\ &= \langle \phi^{\sigma_1} | \otimes \dots \langle \phi^{\sigma_N} | \hat{O} | \phi^{\sigma_N} \rangle \otimes \dots \otimes | \phi^{\sigma_1} \rangle \end{aligned}$$

overlay now yields nested pairs

Linear combinations:

$$\langle \psi^\beta | = A^{\dagger\beta}_{\sigma_N \dots \sigma_1} \langle \phi^{\sigma_1 \dots \sigma_N} | = A^{\dagger\beta}_{\bar{\sigma}} \langle \bar{\phi}^{\bar{\sigma}} |$$

reversed index order

Complex conjugation:

$$A^{\dagger\beta}_{\bar{\sigma}} = \overline{A^{\sigma}_{\beta}}$$

Orthonormality

If $|\phi_\sigma\rangle$ form orthonormal basis: $\langle \phi^\sigma | \phi_\sigma \rangle = \delta^\sigma_\sigma$

If $|\psi_\alpha\rangle$ form orthonormal basis, too: $\langle \phi^\alpha | \psi_{\alpha'} \rangle = \delta^\alpha_{\alpha'}$

Combined:

$$\begin{aligned} \delta^\alpha_{\alpha'} &= \langle \psi^\alpha | \psi_{\alpha'} \rangle = A^{\dagger\alpha}_{\sigma} \underbrace{\langle \sigma | \sigma' \rangle}_{\delta^\sigma_{\sigma'}} A^{\sigma'}_{\alpha'} \\ &= A^{\dagger\alpha}_{\sigma} A^{\sigma}_{\alpha'} = (A^\dagger A)^\alpha_{\alpha'} \end{aligned}$$

Hence A is unitary:

$$I = A^\dagger A \Rightarrow A^\dagger = A^{-1}$$

Operators

$$\hat{O} = |\phi_{\bar{\sigma}}\rangle \hat{O}^{\bar{\sigma}}_{\bar{\sigma}'} \langle \bar{\phi}^{\bar{\sigma}'}|$$

$$O^{\bar{\sigma}}_{\bar{\sigma}'} = \langle \bar{\sigma} | \hat{O} | \bar{\sigma}' \rangle$$

$$O \begin{array}{c} \uparrow \bar{\sigma}' \\ \bullet \\ \downarrow \bar{\sigma} \end{array}$$

bra indices:

A or A^\dagger : w/ indices
incoming arrows

ket indices

A or A^\dagger : down
outgoing arrows

Simplified notation

It is customary to simplify notational conventions for kets and bras.

In kets, use subscript indices as ket names: $|\phi_{\vec{\sigma}}\rangle \equiv |\vec{\sigma}\rangle \equiv |\sigma_1, \sigma_2, \dots, \sigma_N\rangle$
 $\equiv |\sigma_N\rangle \otimes \dots \otimes |\sigma_1\rangle$

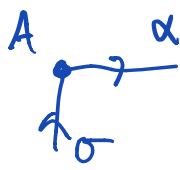
In bras, use superscript indices as bra names:

$$\langle \phi_{\vec{\sigma}} | \equiv \langle \vec{\sigma} | \equiv \langle \sigma_1, \sigma_2, \dots, \sigma_N | \equiv \langle \sigma_1 | \otimes \langle \sigma_2 | \otimes \dots \otimes \langle \sigma_N |$$

Now up/down convention for indices is no longer displayed but it still implicit!

Linear combination of kets: $|\alpha\rangle = |\sigma\rangle A_{\alpha}^{\sigma}$

bra index
↓
↑
ket index

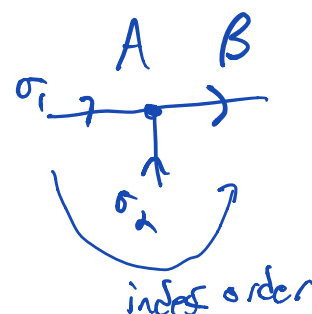


Coefficient matrix = overlap:

$$A_{\alpha}^{\sigma} = \langle \sigma | \alpha \rangle$$

If direct products are involved:

$$|\beta\rangle = |\sigma_2\rangle \otimes |\sigma_1\rangle A_{\beta}^{\sigma_1 \sigma_2}$$



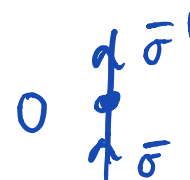
Coefficient matrix = overlap:

$$A_{\beta}^{\sigma_1 \sigma_2} = \langle \sigma_1 | \otimes \langle \sigma_2 | \beta \rangle$$

Operators:

$$\hat{O} = |\vec{\sigma}\rangle \hat{O}_{\vec{\sigma}'}^{\vec{\sigma}} \langle \vec{\sigma}'|$$

$$\hat{O}_{\vec{\sigma}'}^{\vec{\sigma}} = \langle \vec{\sigma} | \hat{O} | \vec{\sigma}' \rangle$$

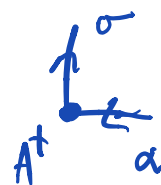


In the overlaps:

bra indices: upper on A or A^\dagger , as incoming arrows

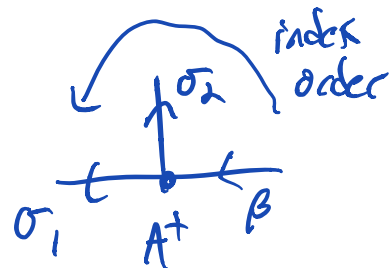
ket indices: lower on A or A^\dagger , as outgoing arrows

Linear combo of bras: $\langle \alpha | = A^{\dagger\alpha}_{\sigma} \langle \sigma |$



Coefficient matrices: $A^{\dagger\alpha}_{\sigma} = \langle \alpha | \sigma \rangle = \overline{\langle \sigma | \alpha \rangle} = \overline{A^{\sigma}_{\alpha}}$ (Hermitian conjugate)

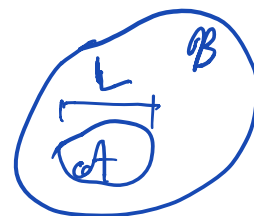
Direct products: $\langle \beta | = A^{\dagger\beta}_{\sigma_2 \sigma_1} \langle \sigma_1 | \otimes \langle \sigma_2 |$



Coefficient matrices: $A^{\dagger\beta}_{\sigma_2 \sigma_1} = \langle \beta | \sigma_2 \rangle \otimes \langle \sigma_1 |$
 $= \overline{\langle \sigma_1 | \otimes \langle \sigma_2 | \beta \rangle} = \overline{A^{\sigma_1 \sigma_2}_{\beta}}$

4. Entanglement entropy and Area Laws (introductory comments)

Consider a quantum system in state $|\psi\rangle$, with density matrix $\hat{\rho} = |\psi\rangle\langle\psi|$



Divide system into two parts, A, B . Suppose A has linear dimension L .

To obtain reduced density matrix of A (or B), trace out B (or A)

'reduced density matrix' for A : $\hat{\rho}_A \equiv \text{Tr}_B \hat{\rho}$ ($\hat{\rho}_B \equiv \text{Tr}_A \hat{\rho}$)

'Entanglement entropy' of A and B :

$$S_{A/B} = -\text{Tr}_A \hat{\rho}_A \log_2 \hat{\rho}_A = -\sum_{\alpha} w_{\alpha} \log_2 w_{\alpha}$$

\uparrow
eigenvalues of ρ_A

Key result: for Hamiltonians with only local interactions, $S_{A/B}$ is governed by an 'area law':

$$S \equiv S_{A/B} \sim (\text{area of boundary of } A) \equiv \partial A$$

$$\sim L^2$$

in 3D for gapped system

$$\sim L$$

in 2D for gapped system

$$\sim \text{const}$$

in 1D for gapped system

$$\sim \text{const} + \ln L$$

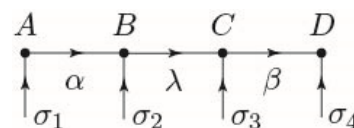
in 1D for *gapless system
 $\overline{\text{I}}_A \text{ TOL, } E_1 - E_0 > 0$

Excited states?

Now consider an MPS of maximal bond dimension D :

$$|4\rangle = |\sigma_4\rangle \otimes |\sigma_3\rangle \otimes |\sigma_2\rangle \otimes |\sigma_1\rangle A_{\alpha}^{\sigma_1} B_{\lambda}^{\sigma_2} C_{\beta}^{\sigma_3} D^{\sigma_4}$$

$$\equiv I^{\lambda \bar{\lambda}} C_{\bar{\lambda} \beta}^{\sigma_3} D^{\sigma_4}$$



Divide system into two parts: Left \rightarrow 2 sites, Right \rightarrow 2 sites

A B

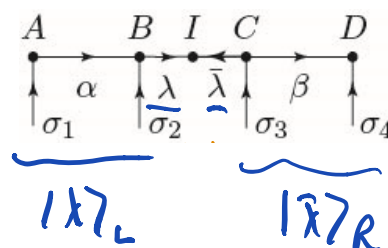
$$|4\rangle = I^{\lambda \bar{\lambda}} |\sigma_4\rangle \otimes |\sigma_3\rangle C_{\bar{\lambda} \beta}^{\sigma_3} D^{\sigma_4}$$

$$\otimes |\sigma_2\rangle \otimes |\sigma_1\rangle A_{\alpha}^{\sigma_1} B_{\lambda}^{\sigma_2}$$

$$= \sum_{\lambda=1}^D I^{\lambda \bar{\lambda}} |\bar{\lambda}\rangle_R \otimes |\lambda\rangle_L$$

implicit sum over $\bar{\lambda}$

$$= \sum_{\lambda=1}^D |\lambda\rangle_R \otimes |\lambda\rangle_L$$



= entangled superposition of two state spaces,
each having dimension of at most D

(After the sum over $\bar{\lambda}$ has been performed explicitly using Kronecker delta, result contains non-covariantly paired indices

Density matrix:

$$\hat{\rho} = |4\rangle\langle 4| = \sum_{\lambda \lambda'} |\lambda\rangle_R |\lambda\rangle_L \langle \lambda' |_L \langle \lambda' |_R$$

Reduced density matrix:

$$\hat{\rho}_A = \text{Tr}_B \hat{\rho} = \sum_{\mu} \langle \mu |_R \sum_{\lambda \lambda'} |\lambda\rangle_R |\lambda\rangle_L \langle \lambda' |_L \langle \lambda' |_R |\mu\rangle_R$$

$$= \sum_{\lambda \lambda'} |\lambda\rangle_L (\rho_A)_{\lambda \lambda'} \langle \lambda' |_L$$

Complete
set of states
 $I_B = \sum_{\mu} |\mu\rangle_R \langle \mu |_R$

with matrix elements

$$(\rho_A)_{\lambda'\lambda} = \sum_{\mu} \underbrace{\langle \mu | \lambda \rangle}_I \underbrace{\langle \lambda' | \mu \rangle}_I = \sum_{\mu} \underbrace{\langle \lambda' | \mu \rangle \langle \mu | \lambda \rangle}_I$$

$$= \langle \lambda' | \lambda \rangle = D \times D \text{ matrix}$$

This matrix has rank $\leq D$ (say D)Let w_α be its eigenvalues, with $\alpha = 1, \dots, D$

and normalization

$$1 = \text{Tr } \hat{\rho}_A = \sum_{\alpha} w_{\alpha}$$

Entanglement entropy:

$$S = - \sum_{\alpha=1}^D w_{\alpha} \log_2 w_{\alpha}$$

Maximal if $w_{\alpha} = \frac{1}{D}$ for all α :

$$S \leq - \sum_{\alpha=1}^D \frac{1}{D} \log_2 \frac{1}{D} = \log_2 D$$

1D gapped:

$$D \sim 2^{\text{const}} \quad (\text{independent of system size})$$

1D critical:

$$D \sim 2^{\text{const} \cdot L} \sim \text{power law in } L \rightarrow \text{efficient}$$

2D gapped:

$$D \sim 2^L$$

} not efficient

3D gapped:

$$D \sim 2^{L^3}$$

Conclusion: MPS can encode ground state efficiently for gapped and gapless systems in 1D, but not 2D or 3D!

5. Tensor network diagrams

[Orus 2014, Sec 4.1]

'tensor' = multi-dimensional array of numbers

'rank of tensor' = number of indices = # of legs

rank-0: scalar

$$A \quad \cdot \quad A^\dagger \quad \cdot$$

rank-1: vector

$$\underbrace{A^\sigma}_{\text{dual vector}} \quad \uparrow \sigma \quad A^\dagger_\sigma \quad \downarrow \sigma$$

rank-2: matrix

$$A^\sigma_\alpha \quad \begin{array}{c} \uparrow \sigma \\ \rightarrow \alpha \end{array} \quad A^{\dagger\alpha}_\sigma \quad \begin{array}{c} \downarrow \sigma \\ \leftarrow \alpha \end{array}$$

rank-3: tensor

$$A^{\alpha\sigma}_\beta \quad \begin{array}{c} \uparrow \sigma \\ \rightarrow \alpha \\ \downarrow \beta \end{array} \quad A^{\dagger\beta}_{\alpha\sigma} \quad \begin{array}{c} \downarrow \sigma \\ \leftarrow \alpha \\ \uparrow \beta \end{array}$$

Index contraction: summation over repeated index

$$C^\alpha_\gamma = \sum_{\beta=1}^{D_\beta} A^\alpha_\beta B^\beta_\gamma \equiv A^\alpha_\beta B^\beta_\gamma$$

$$D_\beta = \text{'bond dimension' of index } \beta$$

$$\begin{array}{c} C \\ \alpha \rightarrow \bullet \rightarrow \gamma \end{array} = \begin{array}{c} A \quad B \\ \alpha \rightarrow \bullet \rightarrow \beta \rightarrow \bullet \rightarrow \gamma \end{array}$$

graphical representation of matrix product

'open index' = non-contracted index (here α, γ)

'tensor network' = set of tensors with some or all indices contracted according to some pattern

Examples:

Your try:

$$C = A_{\alpha} B^{\alpha}$$

$$D^{\alpha}_{\beta} = A^{\delta}_{\gamma} B^{\gamma\alpha}_{\mu} \underbrace{\underbrace{\mu\delta}_{\beta}}_{\delta\beta}$$

$$\underset{\bullet}{C} = \underset{\bullet}{A} \xrightarrow{\alpha} \underset{\bullet}{B}$$

$$\xrightarrow{\alpha} \underset{\bullet}{D} \xrightarrow{\beta} = \begin{array}{c} A \\ \gamma \swarrow \searrow \delta \\ B \quad \mu \quad C \\ \alpha \uparrow \downarrow \beta \end{array}$$

$$\underset{\bullet}{E} = \begin{array}{c} A \\ \gamma \swarrow \searrow \delta \\ B \quad \mu \quad C \\ \alpha \uparrow \downarrow \end{array}$$

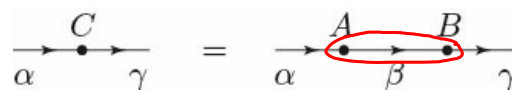
$$\underset{\bullet}{T} = \begin{array}{cc} A & B \\ \delta \uparrow \alpha & \downarrow \beta \\ D & C \\ \gamma & \end{array}$$

Answer:

Cost of computing contractions

Result of contraction does not depend on order in which indices are summed, but numerical cost does!

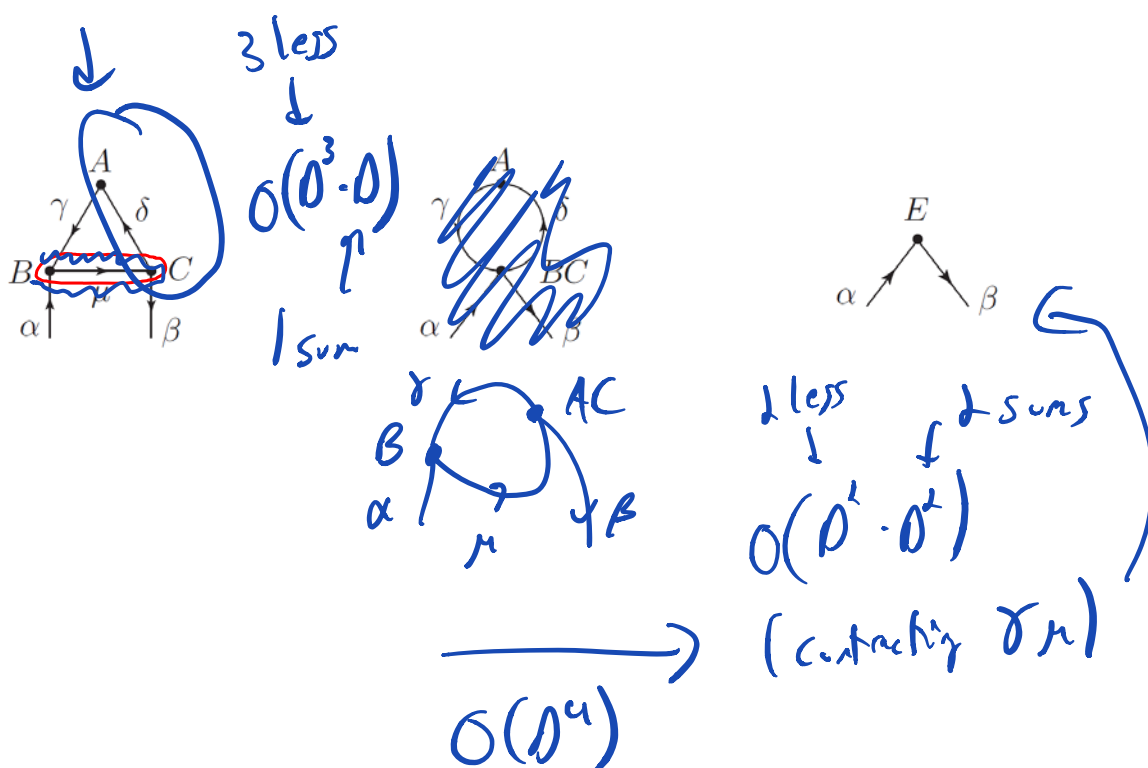
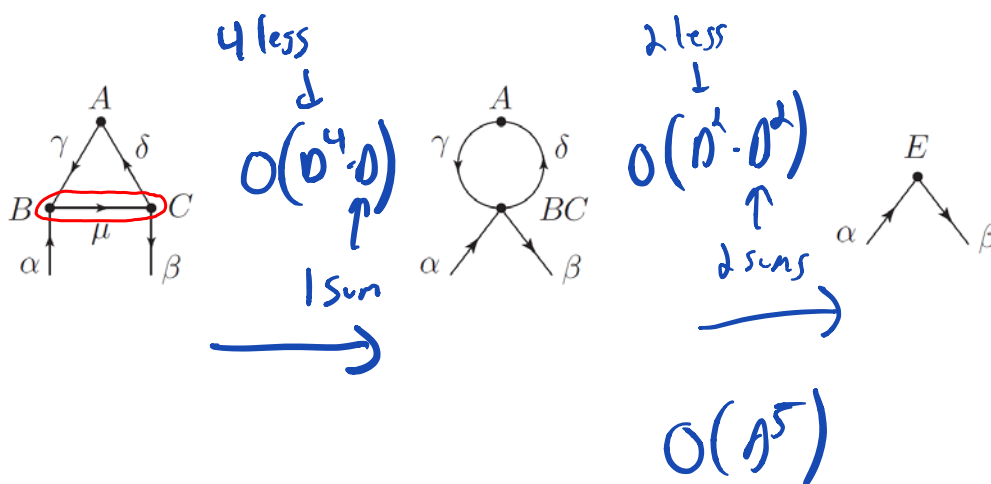
Example 1: cost of matrix multiplication is $O(D^3)$



For every fixed α and γ ($D_\alpha \times D_\gamma$ combinations), sum over D_β values of β

Cost = $D_\alpha \cdot D_\beta \cdot D_\gamma$ (simplifies to D^3 if all bond dimensions are D)

Example 2:



First contraction scheme has total cost $O(D^5)$, second has $O(D^4)$!

Finding optimal contraction order is difficult. In practice, trial and error....

NP complete

In first part of course we will focus on 1D tensor networks, then 2D.

6. Singular value decomposition (SVD)

[Schollwoeck2011, Sec 4]

Any matrix M of dimension $D \times D'$ can be written as $M = U S V^\dagger$

$$\begin{array}{l}
 \underline{D \leq D'}: \\
 \begin{array}{c} D \quad D' \\ \boxed{} \\ M \end{array} = \begin{array}{c} D \quad D \quad D' \\ \boxed{} \quad \boxed{} \quad \boxed{} \\ U \quad S \quad V^\dagger \end{array} \\
 \\
 \underline{D \geq D'}: \\
 \begin{array}{c} D' \\ D \quad \boxed{} \\ M \end{array} = \begin{array}{c} D' \quad D' \quad D' \\ \boxed{} \quad \boxed{} \quad \boxed{} \\ U \quad S \quad V^\dagger \end{array}
 \end{array}$$

Properties of S

- square matrix of dimension $D_{\min} \times D_{\min}$ ($D_{\min} = \min(D, D')$)
- diagonal with non-negative diagonal elements (singular values) $S_{\alpha} \equiv S_{\alpha\alpha}$
- Schmidt rank r : number of non-zero singular values
- arrange in descending order: $s_1 \geq s_2 \geq \dots \geq s_r > 0$

$$S = \text{diag}(s_1, \dots, s_r, \underbrace{0, \dots, 0}_{D_{\min} - r})$$

Properties of U:

- matrix of dimension

$$D \leq D': \quad \begin{matrix} U^\dagger \\ \text{4 rows} \end{matrix} \quad \begin{matrix} U \\ \text{4 columns} \end{matrix} = \begin{matrix} 1 \\ \text{4x4 square with diagonal} \end{matrix}$$

$D_{\min} = D$

- columns are orthonormal

$$U^\dagger U = I$$

Properties of V^\dagger :

- matrix of dimension

$$D \leq D': \quad \begin{matrix} V^\dagger \\ \text{4 rows} \end{matrix} \quad \begin{matrix} V \\ \text{4 columns} \end{matrix} = \begin{matrix} 1 \\ \text{4x4 square with diagonal} \end{matrix}$$

- rows are orthonormal

$$V^\dagger V = I$$

$$D \geq D': \quad \begin{matrix} V^\dagger \\ \text{4 rows} \end{matrix} \quad \begin{matrix} V \\ \text{4 columns} \end{matrix} = \begin{matrix} 1 \\ \text{4x4 square with diagonal} \end{matrix}$$

Useful properties:

$$M M^\dagger = U S V^\dagger \cdot V S U^\dagger = U S^\dagger U^\dagger$$

$$\Rightarrow M M^\dagger U = U S^\dagger$$

$$M^\dagger M = V S U^\dagger U S V^\dagger = V S^\dagger V^\dagger$$

$$M^\dagger M V = V S^\dagger$$

So columns of U are eigenvectors of $M M^T$

Columns of V are eigenvectors of $M^T M$

Truncation

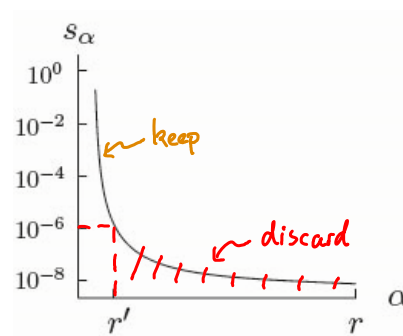
SVD yields optimal approximation of rank r matrix M by a rank r matrix M'

(optimal wrt Frobenius norm: _____)

Suppose $M = U S V^T$

with

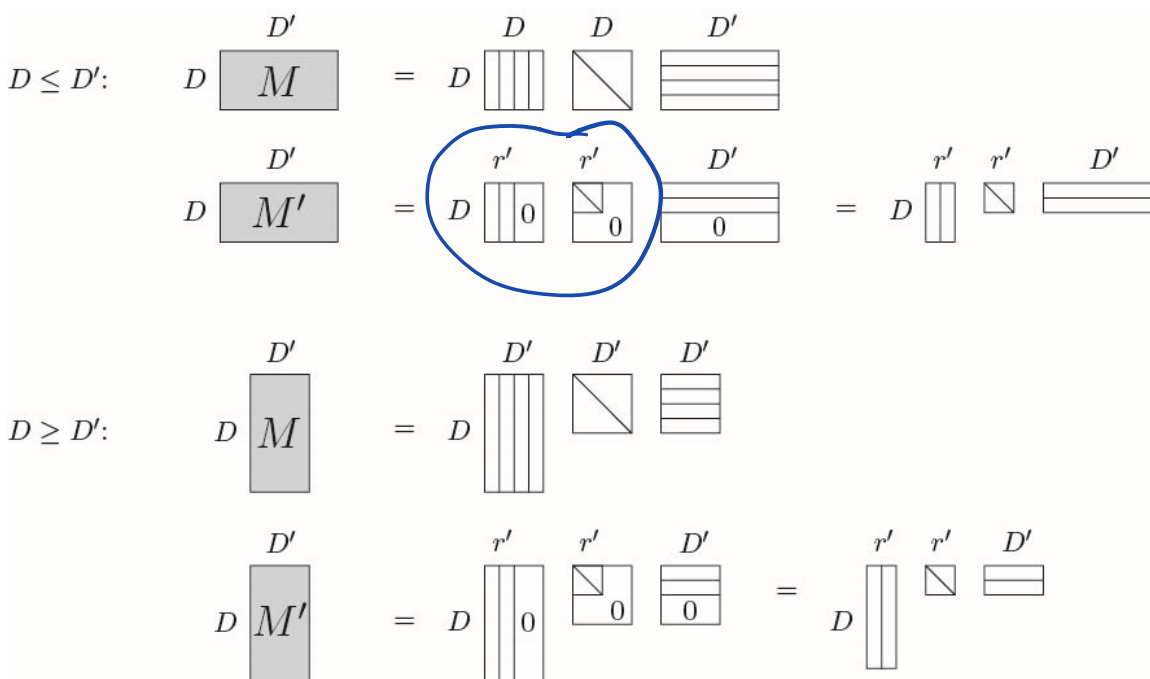
Truncate: $M' = U S' V^T$



with $S' = \text{diag}(s_1, \dots, s_{r'}, \underbrace{0, \dots, 0}_{D_{\min} - r' \text{ zeros}})$

Retain only r' largest singular values!

Visualization, with $r = \text{dim}$:



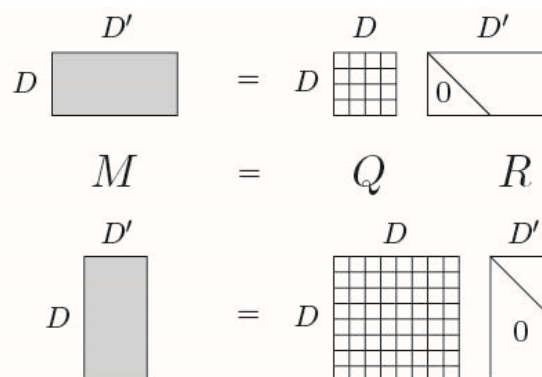
QR decomposition

If singular values are not needed,

a $D \times D'$ matrix M

has the 'full QR decomposition'

$$M = QR$$



with Q a $D \times D$ unitary matrix

$$QQ^T = Q^T Q = I$$

and R a $D \times D'$ upper triangular matrix

$$R_{\alpha\beta} = 0 \text{ if } \alpha > \beta.$$

If $D \geq D'$, then M has the 'thin QR decomposition'

$$M = (Q_1, Q_2) \cdot \begin{pmatrix} R_1 \\ 0 \end{pmatrix} = Q_1 \cdot R_1$$

with $\dim(Q_1) = D \times D'$, $\dim(R_1) = D' \times D'$,

$$\begin{array}{|c|c|} \hline Q_1 & Q_2 \\ \hline \end{array} \cdot \begin{array}{|c|} \hline R_1 \\ \hline 0 \\ \hline \end{array} = \underbrace{\begin{array}{|c|} \hline Q_1 \\ \hline \end{array}}_{\text{rectangular}} \cdot \underbrace{\begin{array}{|c|} \hline R_1 \\ \hline \end{array}}_{\text{square}}$$

and R_1 upper triangular.

$$Q_1^T Q_1 = I \text{ but } Q_1 Q_1^T \neq I$$

QR is numerically cheaper than SVD but has less information (does not provide rank).